до измены

Spy vs. Spy: Rumor Source Obfuscation

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Political activism

Some people have important, sensitive things to say.



Personal confessions

I think Im schizophrenie... I see and hear things and I have a voice in my head, when I go to sleep it's like sleeping in a busy restaurant!

whisper

Others have less important, but sensitive things to say.



secret whisper

















centralized networks are not truly anonymous!

Compromises in anonymity



anonymity loss extends beyond the network

Distributed messaging



Distributed messaging



Distributed messaging



what can an adversary do?









the adversary can figure out who got the message



• G is the graph representing the social network





Alice passes the message to her neighbors



her neighbors pass the message to theirs



the message spreads in all directions at the same rate



the message spreads in all directions at the same rate



the message spreads in all directions at the same rate



this spreading model is known as the diffusion model



can the adversary locate the message author?

Concentration around the center



the message author is in the "center" with high probability

Rumor source identification



high likelihood

low likelihood

diffusion does not provide anonymity

[Shah, Zaman 2011]

Our goal



• N_T : expected number of nodes with the message at time T












Main result: adaptive diffusion $\rho q \rho q$





provides provable anonymity guarantees!

Line graphs



Iet's start with line graphs



T = 0

• the message author starts a rumor at T=0



T = 1

• with probability α , the left (right) node receives the message



T = 1

the node to the right of the author receives the message



T = 2

the rumor propagates in both directions at the same rate



T = 2

the rumor propagates in both directions at the same rate



T = 3

• α is independent of time or hop distance to message author



T = 3

diffusion on a line is equivalent to two independent random walks

Adversary's observation



N = 5

nodes with the message

can the adversary locate the message author?

Maximum likelihood detection



the node in the middle is the mostly likely author

Maximum likelihood detection





consider a line graph



T = 0

• node 0 starts a rumor at T = 0



T = 1

• with probability 1/2, the left (right) node receives the message



T = 1

right node 1 receives the message



T = 2





T = 2

right node 2 receives the message



T = 3





T = 3

Ieft node 1 receives the message

Adversary's observation



 $N_T = 4$ nodes with the message

can the adversary locate the message author?

Maximum likelihood detection



Maximum likelihood detection 1 2 1 0 2 Likelihoods adaptive diffusion diffusion k N_T $N_T/2$ Probability of detection \approx N_T

d-regular trees



what about d-regular trees?

d-regular trees: diffusion



Iikelihoods concentrate heavily around the "center"





initially, the source is also the "virtual source"



• at T = 1, the author selects one neighbor at random



the author passes h = 1and T = 2 to the chosen neighbor

• at T = 1, the author selects one neighbor at random



the chosen neighbor becomes the new virtual source



• at T = 2, the virtual source passes the message to all her neighbors



• as T transitions from even to odd, the virtual source has two options:

keeping the virtual source token

passing the virtual source token

Symmetry properties



the graph is always symmetric around the virtual source

Keeping the virtual source token



• virtual source token is kept with probability $\alpha = \frac{(d-1)^{\frac{1}{2}-h-1}-1}{T}$



Keeping the virtual source token



all leaf nodes with the message pass it to their neighbors
Passing the virtual source token



current virtual source selects one of its neighbors at random

Passing the virtual source token



previous virtual source passes h = 2 and T = 4 to new virtual source

Passing the virtual source token



new virtual source passes the message to its neighbors which in turn pass it to their neighbors

Adversary's observation



can the adversary locate the message author?

Maximum likelihood detection



high likelihood

low likelihood

all nodes except for the final virtual source are equally likely

Main Theorem

- 1. We spread fast: $N_T \approx (d-1)^{\frac{1}{2}}$
- 2. All nodes except for the final virtual source are equally likely to be the source, hence

$$P(\hat{v}_{ML} = v^*) = \frac{1}{N_T - 1}$$

3. The expected distance between the estimated and true source is at least $\frac{T}{2}$.

General graphs



adaptive diffusion for general graphs?

Simulation results: Facebook graph



likelihoods can be approximated numerically

Extensions and related work

Theoretical

- Adversaries with timing information
- Peer-to-peer dynamic networks
- Hiding relays
- Multiple message sources

Systems

- Cyber-bullying detection
- Anonymous video sharing
- Message caching
- Bootstrapping contacts